Emergent Peer-to-Peer Multi-Hub Topology

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Federated Learning [1]



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Algorithm 1 FederatedAveraging. The K clients are indexed by k; B is the local minibatch size, E is the number of local epochs, and η is the learning rate.

Server executes:

initialize w_0 for each round t = 1, 2, ... do $m \leftarrow \max(C \cdot K, 1)$ $S_t \leftarrow$ (random set of m clients) for each client $k \in S_t$ in parallel do $w_{t+1}^k \leftarrow$ ClientUpdate (k, w_t) $m_t \leftarrow \sum_{k \in S_t} n_k$ $w_{t+1} \leftarrow \sum_{k \in S_t} \frac{n_k}{m_t} w_{t+1}^k$ // Erratum⁴

ClientUpdate(k, w): // Run on client k $\mathcal{B} \leftarrow (\text{split } \mathcal{P}_k \text{ into batches of size } B)$ for each local epoch i from 1 to E do for batch $b \in \mathcal{B}$ do $w \leftarrow w - \eta \nabla \ell(w; b)$ return w to server

¹https://en.wikipedia.org/wiki/Federated_learning

Gossip Learning [2]



Algorithm 1	Gossip	Learning	Scheme
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1: initModel()

2: **loop**

- $wait(\Delta)$ 3:
- $p \leftarrow \text{selectPeer()}$ 4:
- send modelCache.freshest() to p5:
- 6: end loop7: procedure ONRECEIVEMODEL(m)
- modelCache.add(createModel(*m*, lastModel)) 8:
- lastModel $\leftarrow m$ 9:

10: end procedure

Decentralized peer sampling [3]

PROOFS Algorithm (active thread)

initial peer list: cache cache size: c shuffle length: 1 node address: **p** 1 **loop** 2 wait(Δ) subset \leftarrow selectRandomSubset(cache, l) 3 $q \leftarrow selectRandom(subset)$ 4 subset.remove(q) 5 subset.add(p) 6 send(q, subset) 7 $subset_q \leftarrow receive(q)$ 8 9 subset_{*a*}.remove(p) subset[°]_q.removeAll(cache) 10 cache \leftarrow subset_a 11





Before and after a shuffling operation. Node 1 sends addresses {itself, 2, 3} to node 4. Node 4 sends back {5,6, 8}.

Random-graph Topology [4]



Power-law Topology [5]



Hub-based topology & Hub sampling

- The objectif is to obtains an overlay network with **h** defined hubs, with **h** a parameter of the algorithm, and each hub is connected to all the nodes in the networks. The application running on top will be able to take advantage of this overlay to speed up message transmission in the network.
- **Preferential Attachment**: Drawing from the concept pioneered by Barabási and Albert [5], preferential attachment dictates that new connections in the network are established preferentially with nodes possessing a higher number of existing connections. This mechanism enables the organic emergence of hubs within the network, with selected nodes naturally assuming central roles based on their connectivity without any explicit distinction other than their number of incoming links.
- Random Attachment: Inspired by gossip-based peer sampling algorithms ([3], [6]), random attachment ensures that nodes maintain connections with a representative and diverse subset of the network. This strategy promotes network robustness by preventing excessive clustering and dependency on specific nodes (hubs). When existing hubs disappear (e.g., due to failures or departure), other nodes within the network are opportunistically elevated to hub status, ensuring continuity and adaptability of the network topology over time.

Algorithm

2

5

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Elevator Algorithm (active thread)

```
initial peer list: cache
   cache size: c
   desired number of hubs: h
  initial backward list: backward_peers (empty)
 1 loop
     wait(\Delta)
     frequency map \leftarrow {}
 3
     for peer in cache do
 4
        peer cache \leftarrow send(CACHE REQUEST, peer)
       frequency_map \leftarrow frequency_map \cup peer_cache
     preferred \leftarrow frequency_map.sort().select(c)
     preferred backward \leftarrow {}
     for peer in preferred do
        peer backward peers \leftarrow send(BACKWARD REQUEST, peer)
       preferred_backward ← preferred_backward ∪ peer_backward_peers
     cache \leftarrow {}
12
     cache \leftarrow selectRandom(preferred, h) + selectRandom(peer_backward_peers, c - h)
13
```

Simulation

• Simulations done with the Java PeerSim [7] simulator (modified),

with the cycle based mode

• Comparaisons against 3 peer sampling algorithms : Newscast, Proofs and Phenix (Power-law)

	Value
Network size	
Number of cycles for each simulation	
Number of times each simulation was run (with different seed)	
c parameter (cache size)	
h parameter (number of hubs)	

- All simulations were run on 16 vCPU, using 64G of memory.
- Code is available at https://gitlab.lip6.fr/legheraba/elevator

Distribution of indegree



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Metrics









Contexts









Conclusion & next steps

Our algorithm allow the emergence of hubs and compared with PROOFS, Newscast and Phenix, our algorithm has equivalent results in term of shortest path length and diameter and is resilient against failures and attacks.

Next steps:

- Add machine learning & compare with other decentralized machine learning approaches ([2], [8])
- Adapt the algorithm to real physical networks
- Implement more complex attacks

Thanks

Any questions ?

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Appendix

Appendix

Background thread

Elevator Algorithm (background thread)

loop request, peer ← receive() if request == CACHE_REQUEST then | send(cache, peer) | backward_peers.add(peer) if request == BACKWARD_REQUEST then | send(backward_peers, peer)

Appendix

Phenix algorithm

Phenix Algorithm (active thread)

p	eer)		
	cache size : c		
	initial peer list: cache		
	initial backward list: backward_peers (empty)		
	number of preferential connections: s		
1	loop		
2	$ $ wait(Δ)		
3	$G_{\text{random}}, G_{\text{friend}} \leftarrow \text{split(cache)}$		
4	cache \leftarrow {}		
5	cache.append(G_{random})		
6	$G_{\text{candidates}} \leftarrow \{\}$		
7	for peer in G_{friend} do		
8	$ $ neighbor_list \leftarrow send(peer, CACHE_REQUEST)		
9	$G_{\text{candidates}} \leftarrow G_{\text{candidates}} \cup \text{neighbor_list}$		
10	$\operatorname{sort}(G_{\operatorname{candidates}})$		
11	$G_{\text{preferred}} \leftarrow G_{\text{candidates}}[0(s-1)]$		
12	for peer in $G_{\text{preferred}}$ do		
13	send(peer, CONNEXION_REQUEST)		
14	$ $ cache.append($G_{ m preferred}$)		